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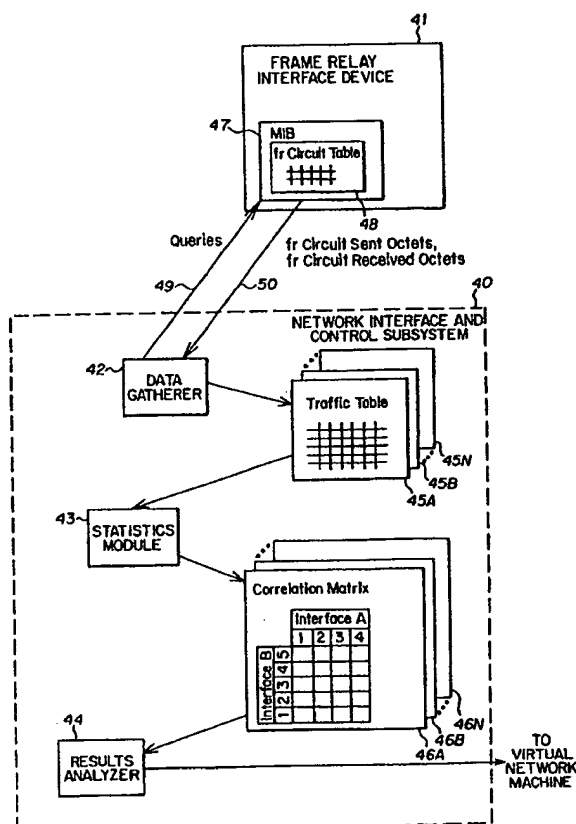
(51) International Patent Classification ⁶ : H04L 12/26		A1	(11) International Publication Number: WO 98/44682
			(43) International Publication Date: 8 October 1998 (08.10.98)
(21) International Application Number: PCT/US98/05370			(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GE, GH, GM, GW, HU, ID, IL, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, ML, MR, NE, SN, TD, TG).
(22) International Filing Date: 18 March 1998 (18.03.98)			
(30) Priority Data: 08/827,541 28 March 1997 (28.03.97) US			
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Published*With international search report.**Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.*

(54) Title: METHOD AND APPARATUS FOR DETERMINING FRAME RELAY CONNECTIONS

(57) Abstract

A system determines that a frame relay connection exists between an interface on a first interface device and an interface on a second interface device, for example by comparing addresses resident in routing tables of the interface devices. The system then queries each of the first and second interface devices, to determine the amount of traffic that is communicated by each channel of each of the two interfaces. This data is then correlated to determine the relative amount of data by each channel on each of the two interfaces, and the pair of channels having the best correlation are determined to represent an actual connection.



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**METHOD AND APPARATUS FOR
DETERMINING FRAME RELAY CONNECTIONS**

BACKGROUND OF THE INVENTION

1. **Field of the Invention**

The invention relates to frame relay communications, and more particularly to a method and apparatus for identifying on which channel a frame relay communication is made.

2. **Discussion of the Related Art**

Computer networks are widely used to provide increased computing power, sharing of resources and communication between users. One type of network is referred to as a frame relay network -- it complies with an interface standard that provides statistical time division multiplexing and circuit switching, and is typically optimized for the transport of protocol-oriented data.

In a typical frame relay network, a communications carrier provides a designated data throughput between two particular entry points (sites) to the network. The data is passed through a "connection," which is also referred to as a "virtual circuit." A permanent virtual circuit (PVC) has one guaranteed data throughput, while a switched virtual circuit (SVC) allows additional throughput in response to a specific demand. Each connection (i.e., virtual circuit) is identified by a data link connection identifier or index (DLCI) number. Many DLCI numbers have local significance only, i.e., a connection identified by a first DLCI number at a first site, may be identified by a different DLCI number at a second site. The connection may be implemented via a set data path, or the data may be re-routed through a different data path depending upon circuit performance. Additional detail on frame relay architectures, protocols, and analysis is described in Analyzing Broadband Networks, Mark A. Miller, M&T Books, pp. 31-169 (1994), which is hereby incorporated by reference.

In the operation and maintenance of computer networks, a number of issues arise including traffic overload on various parts of the network, optimum placement and interconnection of network resources, security, isolation of network faults, and the like. These issues become increasingly complex and difficult to understand and manage as the network becomes larger and more complex. For example, if a network device is not sending messages, it

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may be difficult to determine whether the fault is with the device itself, a data communication link, or an intermediate network device between the sending and receiving devices.

Network management systems have been developed to resolve such issues. Older systems typically operated by collecting large volumes of information, which then required evaluation by a network administrator; this placed a tremendous burden on, and required a highly-skilled network administrator. Newer network management platforms systematize the knowledge of the networking expert such that common problems can be detected, isolated and repaired, either automatically or with the involvement of less-skilled personnel. Such a system typically includes a graphical representation of that portion of the network being monitored by the system. An example of the newer automated system is the Spectrum™ network management platform available from Cabletron Systems, Inc., Rochester, New Hampshire, USA.

It would be desirable to provide a network management system that can determine and display a graphical representation of particular connections which exist in a frame relay network. However, a single interface device, such as a router, switch, or bridge, may have several interfaces, each of which communicates with the frame relay network. Additionally, each interface may communicate over several distinct frame relay connections. Still further, each interface device may not have specific knowledge regarding the other interface device(s) with which it is communicating, much less the specific virtual circuits (connections) that exist at any given time. Accordingly, it would be desirable to be able to determine particular connections from data which is available from the interface devices.

SUMMARY OF THE INVENTION

According to one embodiment of the present invention, a network management system is provided which determines that a frame relay connection is coupled between a first interface on a first interface device and a second interface on a second interface device, for example by comparing addresses which reside in routing tables of the first and second interface devices. The system queries the management information base (MIB) of each of the first and second interface devices, to determine the amount of traffic that is communicated by each channel of the two interfaces. This data is then correlated to determine a relative amount of data transmitted on each channel of the two interfaces, and the two channels having the best correlation are determined to represent a particular connection.

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According to a method embodiment of the invention, a method is provided for identifying a connection in a frame relay network. The method comprises the steps of identifying a first plurality of channels of a first interface and a second plurality of channels of a second interface that potentially relate to a frame relay connection, and then selecting, as an identifier for the frame relay connection, a first channel from the first plurality of channels and a second channel from the second plurality of channels based upon a relative amount of data transferred on each of the first and second channels.

Yet another apparatus embodiment includes a data module and a results analyzer. The data module is coupled to a frame relay network and has an output that provides an indication of a first plurality of channels and a second plurality of channels that potentially relate to a frame relay connection, and the results analyzer selects a first channel from the first plurality of channels and a second channel from the second plurality of channels based upon a relative amount of data transferred on each of the first and second channels, in order to identify the frame relay connection.

In any of these embodiments, a plurality of correlation factors may be determined, for example by a statistics module, for each potential pair of channels. The correlation factor having a best (e.g., highest) value may be used to determine which of the plurality of potential pairs of channels has the best (e.g., highest) correlation factor. A correlation factor may be related to a similarity between an amount of data transmitted by a first channel and an amount of data received by a second channel during a period of time.

In any of these embodiments, a first interface device may be queried to determine an amount of data that has been transferred by one channel as of a first time, and the first interface device queried again to determine an amount of data that has been transferred by the one channel as of a second time, wherein the amount of data being transferred is determined to be the difference between these two amounts. A similar set of queries and difference determination may be made for a second interface of the potential pair.

In accordance with at least one embodiment, a determination may be made that a connection exists between an interface of a first interface device and an interface of a second interface device by comparing an address of the first interface with a destination address of data transmitted by the second interface.

These and other features and benefits of the present invention will be more particularly described in the following detailed description.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a schematic diagram of a network management system that interfaces with elements of a frame relay network;

5 Fig. 2 is a block diagram of two devices which each interface with a frame relay network, such as the one shown in Fig. 1;

Fig. 3 shows data which may be obtained from the two devices of Fig. 2, in accordance with at least one network protocol standard;

10 Fig. 4 is a block diagram of a network management system according to one embodiment of the invention, and a frame relay device which communicates therewith;

Fig. 5 is a process flow diagram showing steps of a process in accordance with an embodiment of the invention;

Fig. 6 is a flow diagram with greater detail of a portion of the process shown in Fig. 5;

15 Fig. 7 is a flow diagram with greater detail of another portion of the process shown in Fig. 5;

Fig. 8 is an example of a traffic table, for data such as that shown in Fig. 3, wherein the data relates to a first interface;

Fig. 9 is an example of a traffic table, for data that relates to a second interface; and

20 Fig. 10 is an example of a correlation matrix for the data that is shown in Fig. 3.

DETAILED DESCRIPTION

Aspects of the present invention will be described in conjunction with the use of a model-based network management system, such as Spectrum™, which is described in U.S. Patent No. 5,261,044, issued November 9, 1993 to R. Dev et al., and hereby incorporated by
25 reference in its entirety. The Spectrum™ network management system is commercially available and described in various user manuals and literature available from Cabletron Systems, Inc., 35 Industrial Way, Rochester, New Hampshire 03867.

Other commercially available network management systems and related applications include: HP OpenView, Hewlett Packard Corp., 3000 Hanover Street, Palo Alto, CA 94304;
30 LattisNet, Bay Networks, 4401 Great American Pkwy., Santa Clara, CA 95054; IBM Netview/6000, IBM Corp., Old Orchard Road, Armonk, NY 10504; and SunNet Manager, SunConnect, 2550 Garcia Ave, Mountain View, CA 94043. Embodiments of the invention

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described herein may be applicable to such other network management systems or portion thereof.

Fig. 1 shows a network management system 10 that monitors a frame relay network 11 and other network(s) 12. The other network(s) 12 may be other frame relay networks, or other types of networks, e.g., ethernet, token ring, FDDI, etc. The frame relay network 11 is illustrated as connecting three different sites 13 that communicate across the frame relay network - site A (13A), site B (13B), and site C (13C). The network management system 10 monitors the frame relay network 11 by communicating with devices within each of the sites 13A-C. Each device within the sites 13A-C is referred to herein as an "interface device."

Fig. 1 shows one embodiment of the network management system 10, including a network interface and control subsystem 16, a virtual network machine 17, and a user interface 18. The virtual network machine 17 is coupled to both the network interface and control subsystem 16 and the user interface 18. The network interface and control subsystem 16 is coupled to each of the frame relay sites 13A-C by a corresponding one of communication paths 14A-C, and is coupled to the other network(s) 12 by communication path 15. The communication paths 14A-C and 15 may be actual connections, or may represent a virtual access to data that relates to a particular site. For example, the network interface and control subsystem 16 may access a single data repository that stores data for all of the sites 13A-C, instead of directly accessing the interface devices on each site.

In one embodiment, the network management system 10 comprises a programmed digital computer, wherein a program is implemented using an object-oriented programming language such as C++, Eiffel, SmallTalk, and Ada. The virtual network machine 17 includes interrelated intelligent models (objects) of network entities and relations between network entities, including a capability for acquiring network data pertaining to a condition of a network entity from the corresponding network entity, and for those entities not capable of being contacted, inferring their status from the status of other communicating entities. The virtual network machine 17 maintains objects which include network data relating to the corresponding network entity and one or more inference handlers for processing the network data, the inference handlers being responsive to changes occurring in the same and/or a different object. The network data can then be transferred to the user interface 18 for supplying the network data to a user.

The user interface 18 may be resident on a computer that is different from the computer on which subsystem 16 and virtual network machine 17 reside. In such an arrangement, the user

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interface 18 may be referred to as a client, while the network interface and control subsystem 16 and the virtual network machine 17 may be referred to as a server. Several instances of the user interface 18 may exist for a single instance of the virtual network machine 17, so that several users may concurrently view different aspects of the networks being monitored by the network management system 10.

In one embodiment of the invention, the network interface and control subsystem 16 obtains data relating to the frame relay network 11 in accordance with the Simple Network Management Protocol (SNMP). SNMP is a standard that defines common messages and databases to support network management. A group known as the Internet Engineering Task Force (IETF) has established several standards in accordance with the SNMP framework, and in particular a standard known as a Management Information Base (MIB). The MIB specifies the data items that certain devices must maintain, other data items that are proprietary or optional, and the operations allowed on such devices. For example, the MIB specifies that Internet Protocol (IP) software must keep a count of all octets (8-bit bytes) that arrive over each network interface. Thus, a network management system is able to transmit a specific query (defined by the MIB) to a network interface device, and receive in response a message providing such count data stored within the network interface device's MIB. A document known as the IETF RFC (Request for Comments) 1315 defines a portion of the MIB for use with managing frame relay protocols.

RFCs such as RFC-1315 are available in both printed and electronic form. Printed copies may be obtained from Network Solutions, Inc., 505 Huntmar Park Drive, Herndon VA 22070, while electronic copies may be accessed on the internet. In particular, an internet user may use an "anonymous" FTP to the host ds.internic.net (residing at 129.20.239.132) and retrieve the desired files from the directory "rfc/".

Fig. 2 illustrates detail of two interface devices, site A router 21, and site B router 22, each of which is coupled to the frame relay network 11. Each of routers 21 and 22 may have several interfaces. In particular, site A router 21 has interfaces If7, If8, and If9, while site B router 22 has interfaces If1, If2, If3, and If4. Each of these interfaces are coupled to the frame relay network 11.

In order for a network management system to effectively monitor and accurately represent the frame relay network 11, it is desirable to determine which interfaces share a connection within the frame relay network. If two devices are performing routing or bridging,

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then the network management system 10 can infer a connection between two of the interfaces by reading the route tables or transparent bridge tables (defined within the MIB) of the interface devices 21 and 22. In the example shown in Figs. 2 and 3, the interface devices 21 and 22 are routers that perform routing functions, and therefore the IP routing tables may be read by the network management system 10. Information relating to the IP routing tables is defined within the MIB, and is available from RFC documents such as MIB-II RFC-1213.

In particular, Fig. 3 shows router information data 30 and routing table entries 31 for the site A router 21. The router information data 30 indicates that device 21 is a frame relay interface, and that information was requested pertaining to interface If9, which has an IP address of 131.124.196.1. The router table information 31 shows that there are three groups of destination addresses (e.g., subnets) for messages to be sent from interface If9, two of which had a next hop IP address of 131.124.196.3, and one of which had a next hop IP address of 131.124.196.2.

Similarly, router information data 32 relates to the site B router 22, and router table entries 33 contain information pertaining to interfaces If2 and If4 of the site B router 22. The router table information 33 shows that there are three groups of destination addresses (e.g., subnets) for messages to be sent from interface If4 with a next hop IP address of 131.124.196.1.

Because messages are to be sent from interface If9 to interface If4 (as indicated by the next hop IP address 131.124.196.2 in table 31 matching the IP address of interface If4), a network management system can infer that a frame relay connection exists between interface If9 and interface If4. This conclusion is further supported by the evidence that messages are to be sent from interface If4 to interface If9, as indicated by several of the next hop IP addresses in router table entries 33 (from interface If4) matching the IP address of interface If9 (131.124.196.1).

Thus, by determining the IP address of a message to be sent from an interface across a frame relay network, and then matching this IP address to the IP address of another interface on the frame relay network, a network management system can determine which two interfaces on a frame relay network share a connection. Similar results may be achieved by querying the transparent bridge tables if the two interface devices are performing bridge functions. The bridge MIB is defined in RFC-1493.

However, each interface may have several "channels", each communicating across a different connection within the frame relay network 11. For example, Fig. 2 shows that interface

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If9 has three channels 211, 212, and 213, interface If8 has two channels 214 and 215, interface If4 has two channels 221 and 222, interface If3 has one channel 223, and interface If2 has three channels 224, 225, and 226. Thus, even though it may be known which two interfaces share a connection within the frame relay network 11, it still may be uncertain which two channels of the two interfaces actually share the connection.

Fig. 4 is a block diagram showing operation of a system to determine the channels that share a frame relay connection. In particular, Fig. 4 depicts a network interface and control subsystem 40 which represents an embodiment of the network management and control subsystem 16 of Fig. 1. The subsystem 40 communicates with a frame relay interface device 41, which as described above, may be a router, switch, network interface card, or any device that communicates across a frame relay network (not shown in Fig. 4). The network interface and control subsystem 40 includes a data gatherer 42, a statistics module 43, and a results analyzer 44. Traffic tables 45A, 45B, ... 45N are coupled between the data gatherer 42 and the statistics module 43, and correlation matrices 46A, 46B, ..., 46N are coupled between the statistics module 43 and the results analyzer 44.

The results analyzer 44 is the portion of the network interface and control module that provides input to the virtual network machine, and the data gatherer 42 interfaces with the interface device 41. As described above, associated with the interface device 41 is a MIB 47. A portion of the MIB 47 is the frame relay circuit table (frCircuitTable) 48, as defined in RFC-1315. The data gatherer 42 communicates with the interface device 41 in the form of queries 49 which request information from the MIB 47, and in response the MIB 47 provides a count of octets sent by a particular channel of an interface of the interface device (frCircuitSent Octets), and a count of octets received by a particular channel of an interface of the interface device (frCircuitReceivedOctets), generally shown as 50 in Fig. 4.

Fig. 5 is a high level process flow diagram of operations performed by an embodiment of the network interface and control subsystem 40. In step 51, the data gatherer 42 determines that a frame relay connection exists between a first interface and a second interface (e.g., interface If9 of site A router 21 and interface If4 of site B router 22). As discussed above, this determination may be made by analyzing the next hop address within a routing table, or by other methods known to those skilled in the art.

In step 52, the data gatherer 42 determines an amount of data that is communicated in each direction by each channel of the first and second interfaces. This step may include

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determining which channels are actually active (which may also be determined by querying the MIB 47), in which case any channels that are not active may be eliminated from consideration. Additionally, it is not necessary to determine the amount of data in both directions, although such analysis may result in a more complete solution as described in more detail below.

5 In one embodiment, the data gatherer 42 stores the data amounts in the appropriate traffic table 45. One of the traffic tables 45A may be for the analysis of a first interface, while another of the traffic tables 45B may be for the analysis of a second interface.

 Once the amount of data has been determined, the statistics module 43 determines the level of correlation between data on each potential channel pair. A potential channel pair is
10 made up of one channel from each interface for which a connection is known to exist. In particular, step 53 indicates that the level of correlation is determined between the amount of data received by each channel of each interface and the amount of data transmitted by each channel of the other interface. In the embodiment shown in Fig. 4, the correlation data is stored in one of the correlation matrices 46A-N (e.g., one for each interface pair). As shown in correlation matrix
15 46A, the matrix includes several rows for interface B, and several columns for interface A. Each row represents one channel on interface B, and each column represents one channel on interface A. Thus, each row-column intersection can store a correlation factor that relates the amount of data traffic on the interface B channel represented by the corresponding row, with the amount of data traffic on the interface A channel represented by the corresponding column. The statistics
20 module 43 fills in the correlation matrix 46A with correlation factors derived from the traffic data stored in the traffic tables 45.

 Once the correlation matrix 46A is populated, the results analyzer 44 evaluates the correlation factors to determine the channel of interface A and the channel of interface B that are communicating via the frame relay connection determined in step 51, more specifically by
25 inferring that a specific connection exists between highly correlated channels (step 54). Once the specific connection is inferred, the results analyzer 44 may update the virtual network machine 17 to represent the connection (step 55).

 Fig. 6 shows more detail of an implementation of step 52 of Fig. 5. In this implementation, the circuit tables of MIB 47 for each of the first and second interfaces are
30 queried to determine the value of a "frame in counter" and a value of a "frame out counter" for each active channel (step 61). As shown in Fig. 4, one example of a "frame in counter" is the MIB variable frCircuitSentOctets, and an example of a "frame out counter" is the MIB variable

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frCircuitReceivedOctets. These variables each count the total number of octets (8-bit bytes) transmitted or received by a specific channel for which the MIB is queried, including network overhead associated with transferring the data. However, other data may be used to determine the amount of traffic or other relevant statistics, for example a variable that tracks the total
 5 number of frames transmitted or received, or a variable that tracks the total number of errors in transmitted or received data.

In step 62, the values of the "frame in counter" and the "frame out counter" are each stored, and a predetermined amount of time is allowed to pass. Once the predetermined amount of time has passed, the circuit tables are again queried to determine a new value of each of the
 10 "frame in counter" and the "frame out counter." Examples of a predetermined amount of time are 30 seconds, or 50 seconds, although highly correlated data may be obtained with even less or significantly more time. The reason that two queries are performed in some embodiments is that the counters of the circuit table typically count the data from the last time that the interface device was initialized. Thus, if the two interface devices were initialized at different times, the
 15 difference between the values from the first query and the values from the second query will account for the different initialization times.

In step 64, each value from the first query (step 61) is subtracted from a corresponding value from the second query (step 63), to determine the amount of data communicated by each channel in the amount of time used as a delay in step 62. The resulting values may be stored in
 20 traffic table 45A, an example of which is described in more detail below.

Fig. 7 shows more detail of an implementation of steps 53 and 54 of Fig. 5. In this implementation, for each potential channel pair, a correlation factor is calculated which represents the similarity of the amount of data transferred/received by one channel compared with the amount of data transferred/received by the other channel (step 71). In one embodiment,
 25 the correlation factor is calculated from the data transferred in both directions, as indicated in Equation (1):

$$(1) \quad CF_{1,2} = 100 - \frac{(\text{abs}(\Delta In_1 - \Delta Out_2) + \text{abs}(\Delta Out_1 - \Delta In_2)) * 100}{\Delta In_1 + \Delta Out_1 + \Delta In_2 + \Delta Out_2}$$

30

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where:

“abs” is an absolute value function;

DeltaOut_i represents a value from column 87 of Fig. 8 or column 97 of Fig. 9; and

DeltaIn_i represents a value from column 88 of Fig. 8 or column 98 of Fig. 9.

5

If the two channels for which the CF is calculated are connected, the absolute value of the difference between DeltaIn₁ and DeltaOut₂ and the absolute value of the difference between DeltaOut₁ and DeltaIn₂ would each ideally be zero, if the frame counters were accessed at precisely the same times. Thus, for ideal data transmission, the correlation factor CF would be equal to 100. However, it is likely that some differences will remain in these access times. The denominator of equation (1) accounts for the total amount of data transferred by the two channels; in this case, channel “1” and channel “2” for which the correlation factor “CF_{1,2}” is calculated. As indicated above, the correlation matrix 46 may then be populated with the correlation factor CF that is calculated for each potential channel pair (step 72).

15

Once the correlation matrix has been populated, then the correlation factors may be compared with one another to determine the potential channel pair that has the highest level of correlation (step 73). Once this comparison is performed, the potential channel pair having the highest level of correlation is selected as the specific frame relay connection that exists between the first interface and the second interface (step 74).

20

Fig. 8 shows a specific example of a traffic table 80, which contains eight columns, each populated with data. Column 81 (“IfIndex”) identifies the interface for which queries have been performed, in this case interface If9. Column 82 (“DLCI”) identifies the channel number within the interface. In this example, there are three active channels within interface If9: channel 31, channel 103, and channel 351.

25

Column 83 (“OutFrames-t0”) stores the “frame out counter” value for each channel that resulted from the first query, while column 84 (“InFrames-t0”) stores the “frame in counter” value for each channel that resulted from the first query. Similarly, column 85 (“OutFrames-t1”) stores the “frame out counter” value for each channel that resulted from the second query, while column 86 (“InFrames-t1”) stores the “frame in counter” value for each channel that resulted from the second query. From these values, the value for column 87 (“DeltaOut”) may be determined by subtracting the value in column 83 from the value in column 85, for each channel. Similarly, the value for column 88 (“DeltaIn”) can be determined by subtracting the value in

30

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column 84 from the value in column 86, for each channel. DeltaOut represents the amount of octets transmitted by the associated channel during the time interval between the two queries, while DeltaIn represents the number of octets received by the associated channel during this same time interval.

5 Fig. 9 is similar to Fig. 8, but represents an example of traffic data compiled for the channels 40 and 56 of interface If4. As in Fig. 8, column 91 identifies the interface, column 92 identifies the particular channel for which data has been gathered, columns 93 and 94 store the counter values from the first query, and columns 95 and 96 store the counter values from the second query. The values for columns 97 and 98 are then determined by subtraction from the
10 earlier columns.

Fig. 10 shows a specific example of a correlation matrix 100 that contains two columns, one for each of the channels of interface If4 shown in Fig. 2. The correlation matrix 100 further contains three rows, one for each of the channels of interface If9 shown in Fig. 2. The correlation factors have been calculated from equation (1) using the data from Figs. 8 and 9. For
15 example, the correlation factor for the potential channel pair: channel 40 of interface If4, and channel 31 of interface If9, has been determined as:

$$(2) \quad CF_{31-40} = 100 - \frac{(\text{abs}(26,663 - 14,137) + \text{abs}(49,246 - 726)) * 100}{26,663 + 49,246 + 14,137 + 726} = 32$$

20

In this example, the highest correlation factor is for the channel pair of channel 56 of interface If4 and channel 31 of interface If9, since 99 is the highest correlation factor value in the correlation table 100. Therefore, in this example, the virtual network machine 17 would be
25 updated to represent that a specific frame relay connection exists between channel 31 and channel 56.

Generally, only a single connection will exist between any two interfaces, so the highest value of the correlation factor is selected. If two correlation factors are very close, however, this may be indicative of two specific channels. Such a situation may be resolved by determining the
30 amount of data throughput supported by each of the two suspected channels as compared with the amount of data actually being transmitted across the frame relay network.

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Although Fig. 4 shows a particular implementation of elements within the network interface and control subsystem 40, those skilled in the art will realize that other implementations are within the scope of the invention. For example, instead of tables or matrices, the appropriate data may be stored in any number of ways, and the functions of the data gatherer 42, the statistics module 43, and the results analyzer 44 may be re-allocated to a different combination of modules.

The embodiments of the network management system described herein may also be implemented in specially-designed hardware such as a state machine or other digital circuitry, or in any of several programming languages on a general purpose computer, or as a programmed general purpose computer itself.

For example, elements of the network management system shown in Fig. 4 may be implemented as software on a floppy disk, compact disk, or hard drive, which controls a computer, for example a general purpose computer such as a workstation, a mainframe or a personal computer, to perform steps of the disclosed processes or to implement equivalents to the disclosed block diagrams. Such a general purpose computer typically includes a central processing unit (CPU) coupled to random access memory (RAM) and program memory via a data bus. The general purpose computer may be connected to the monitored networks in order to interface with other elements, and may provide commands to devices on the network in order to control the network configuration.

Alternatively, the elements of the embodiments described herein may be implemented as special purpose electronic hardware. Additionally, in either a hardware or software embodiment, the functions performed by these different elements may be combined in varying arrangements of hardware and software.

For example, in at least one embodiment of the invention, the portions of the network interface and control shown in Fig. 4 are supplied as a separate software package known as a "Management Module," which may be purchased and integrated with an existing instance of the Spectrum Network Management System.

Having thus described at least one illustrative embodiment of the invention, various modifications and improvements will readily occur to those skilled in the art and are intended to be within the scope of the invention. Accordingly, the foregoing description is by way of example only and is not intended as limiting.

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CLAIMS

1. A method for identifying a frame relay connection in a network, the method comprising the steps of:

5 identifying a first plurality of channels of a first interface and a second plurality of channels of a second interface that relate to a frame relay connection between the first and second interfaces; and

selecting, as an identifier of the frame relay connection, a first channel from the first plurality of channels and a second channel from the second plurality of channels
10 based upon a relative amount of data transferred on each of the first and second channels.

2. The method of claim 1, wherein the selecting step includes:

identifying a plurality of different potential pairs of channels, each potential pair including one of the first plurality of channels and one of the second plurality of
15 channels;

determining a plurality of correlation factors, each correlation factor relating to a respective one of the potential pairs of channels; and

determining which of the correlation factors has a best value with respect to others of the correlation factors. in order to determine which of the potential pairs of channels
20 has a best correlation factor.

3. The method of claim 2, wherein the step of determining a correlation factor includes determining an amount of similarity between an amount of data transmitted by the one of the first plurality of channels and an amount of data received by the one of the second plurality
25 of channels during a period of time.

4. The method of claim 3, wherein the step of determining a correlation factor further includes determining an amount of similarity between an amount of data received by the one of the first plurality of channels and an amount of data transmitted by the one of the second
30 plurality of channels during the period of time.

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5. The method of claim 1, including the step of determining an amount of data transferred by each of the first plurality of channels and each of the second plurality of channels; and

wherein the step of selecting includes determining the relative amount of data
5 based upon the amount of data transferred by each of the first plurality of channels and each of the second plurality of channels.

6. The method of claim 5, wherein the step of determining, for one channel of the first plurality of channels that resides on a first interface device of the frame relay network,
10 includes:

querying the first interface device to determine a first amount of data that has been transferred on the one channel as of a first time;

querying the first interface devices to determine a second amount of data that has been transferred on the one channel as of a second time; and

15 determining a difference between the first and second amounts of data in order to determine the amount of data being transferred by the one channel.

7. The method of claim 1, wherein the step of identifying includes determining that a connection exists between an interface of a first interface device that is coupled to the frame
20 relay network, and an interface of a second interface device that is coupled to the frame relay network;

wherein the interface of the first interface device includes the first plurality of channels, and the interface of the second interface device includes the second plurality of channels.

25

8. The method of claim 7, wherein the step of determining includes comparing an address of the first interface with a destination address of data transmitted by the second interface.

30 9. An apparatus for identifying a frame relay connection in a network, the apparatus comprising:

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means for identifying a first plurality of channels of a first interface and a second plurality of channels of a second interface that relate to a potential frame relay connection between the first and second interfaces; and

5 means for selecting a first channel from the first plurality of channels and a second channel from the second plurality of channels based upon a relative amount of data transferred on each of the first and second channels, to identify the frame relay connection.

10. The apparatus of claim 9, wherein the means for selecting includes:

10 means for identifying a plurality of different potential pairs of channels, each potential pair including one of the first plurality of channels and one of the second plurality of channels;

means for determining a plurality of correlation factors, each correlation factor relating to a respective one of the potential pairs of channels; and

15 means for determining which of the correlation factors has a best value with respect to others of the correlation factors, in order to determine which of the potential pairs of channels has a best correlation factor.

11. The apparatus of claim 10, wherein the means for determining a correlation factor
20 includes means for determining an amount of similarity between an amount of data transmitted by the one of the first plurality of channels and an amount of data received by the one of the second plurality of channels during a period of time.

12. The apparatus of claim 11, wherein the means for determining a correlation factor
25 further includes means for determining an amount of similarity between an amount of data received by the one of the first plurality of channels and an amount of data transmitted by the one of the second plurality of channels during the period of time.

13. The apparatus of claim 9, including means for determining an amount of data
30 transferred by each of the first plurality of channels and each of the second plurality of channels;

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wherein the means for selecting includes means for determining the relative amount of data based upon the amount of data transferred by each of the first plurality of channels and each of the second plurality of channels.

5 14. The apparatus of claim 13, wherein the means for determining, for one channel of the first plurality of channels that resides on a first interface device of the frame relay network, includes:

 means for querying the first interface device to determine a first amount of data that has been transferred on the one channel as of a first time;

10 means for querying the first interface devices to determine a second amount of data that has been transferred on the one channel as of a second time; and

 means for determining a difference between the first and second amounts of data in order to determine the amount of data being transferred by the one channel.

15 15. The apparatus of claim 9, wherein the means for identifying includes means for determining that a connection exists between an interface of a first interface device that is coupled to the frame relay network, and an interface of a second interface device that is coupled to the frame relay network;

 wherein the interface of the first interface device includes the first plurality of
20 channels, and the interface of the second interface device includes the second plurality of channels.

 16. The apparatus of claim 15, wherein the means for determining includes means for comparing an address of the first interface with a destination address of data transmitted by the
25 second interface.

 17. An apparatus for identifying a frame relay connection in a network, the apparatus comprising:

 a data module, coupled to the network, having an output that provides an
30 indication of a first plurality of channels and a second plurality of channels that relate to a potential frame relay connection; and

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5 a results analyzer, having an input coupled to the output of the data module, the results analyzer selecting a first channel from the first plurality of channels and a second channel from the second plurality of channels based upon a relative amount of data transferred by each of the first and second channels, to identify the frame relay connection.

18. The apparatus of claim 17, wherein the indication of the first plurality of channels and the second plurality of channels includes a plurality of different potential pairs of channels, each of the potential pairs including one of the first plurality of channels and one of the second
10 plurality of channels.

19. The apparatus of claim 18, further comprising a statistics module, coupled between the output of the data module and the input of the results analyzer, the statistics module determining a plurality of correlation factors, each correlation factor relating to a respective one
15 of the potential pairs of channels.

20. The apparatus of claim 19, wherein the results analyzer determines which of the plurality of correlation factors has a best value with respect to others of the plurality of correlation factors, to determine which of the plurality of potential pairs of channels has a best
20 correlation factor.

21. The apparatus of claim 19, wherein the statistics module determines a correlation factor between one of the first plurality of channels and one of the second plurality of channels, the correlation factor being proportional to an amount of similarity between an amount of data
25 transmitted by the one of the first plurality of channels and an amount of data received by the one of the second plurality of channels during a period of time.

22. The apparatus of claim 21, wherein the correlation factor is further proportional to an amount of similarity between an amount of data received by the one of the first plurality of
30 channels and an amount of data transmitted by the one of the second plurality of channels during the period of time.

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23. The apparatus of claim 17, wherein the data module determines an amount of data transferred by each of the first plurality of channels and each of the second plurality of channels;

the apparatus further comprising a statistics module, coupled between the data module and the results analyzer, the statistics module determining the relative amount of data based upon the amount of data transferred by each of the first plurality of channels
5 and each of the second plurality of channels.

24. The apparatus of claim 23, wherein the data module determines the amount of data being transferred by one channel of the first plurality of channels based upon a difference
10 between an amount of data that has been transferred by the one channel as of a first time, and an amount of data that has been transferred by the one channel as of a second time.

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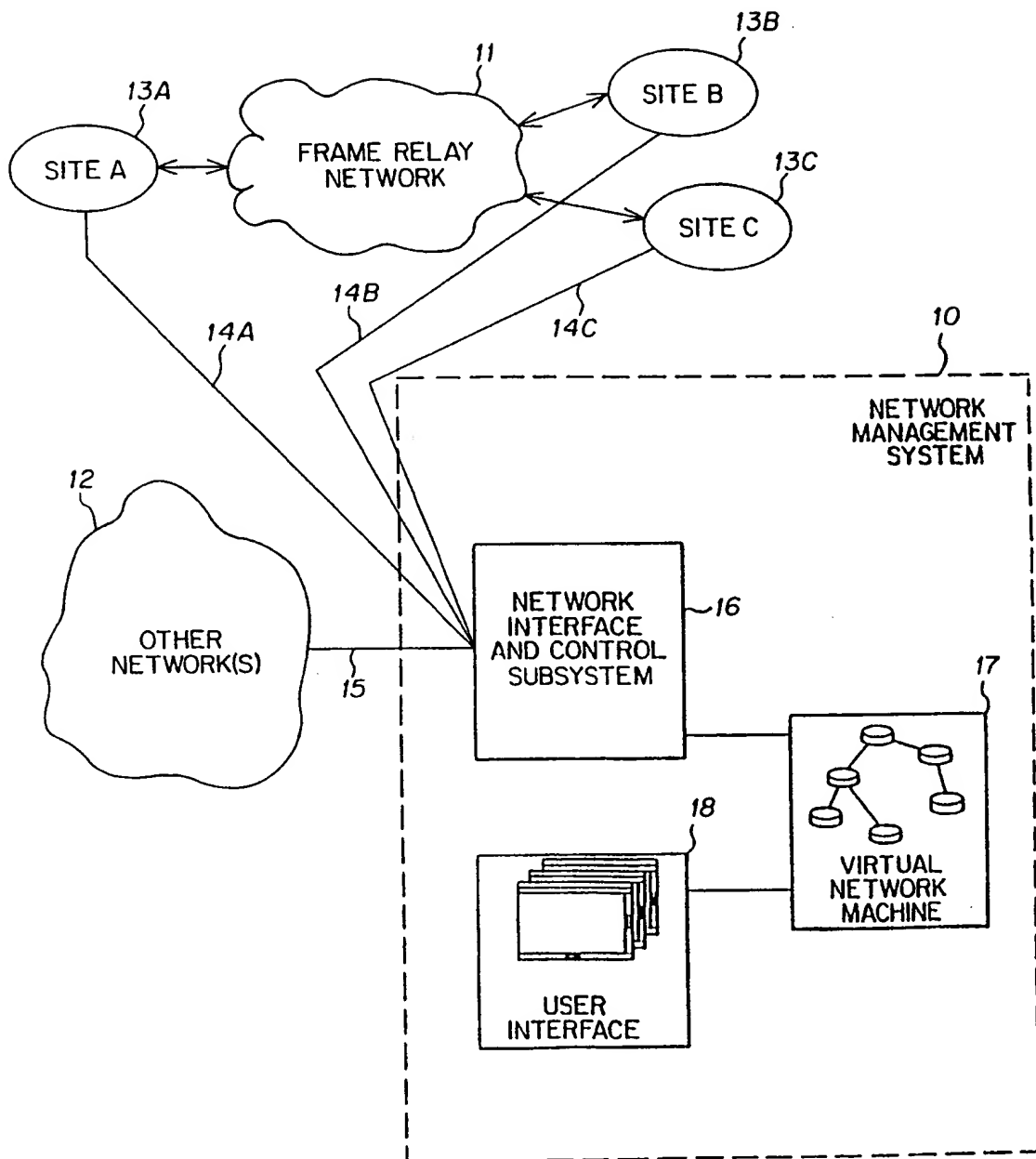


FIG. 1

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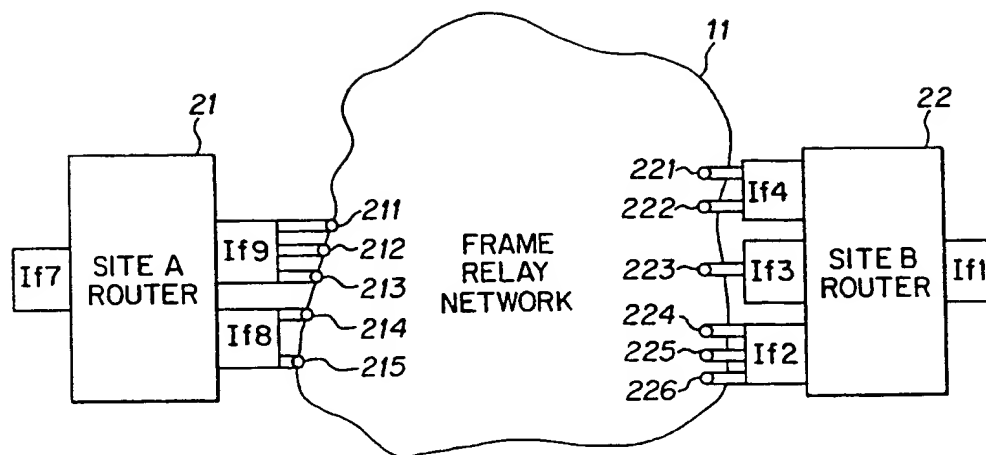


FIG. 2

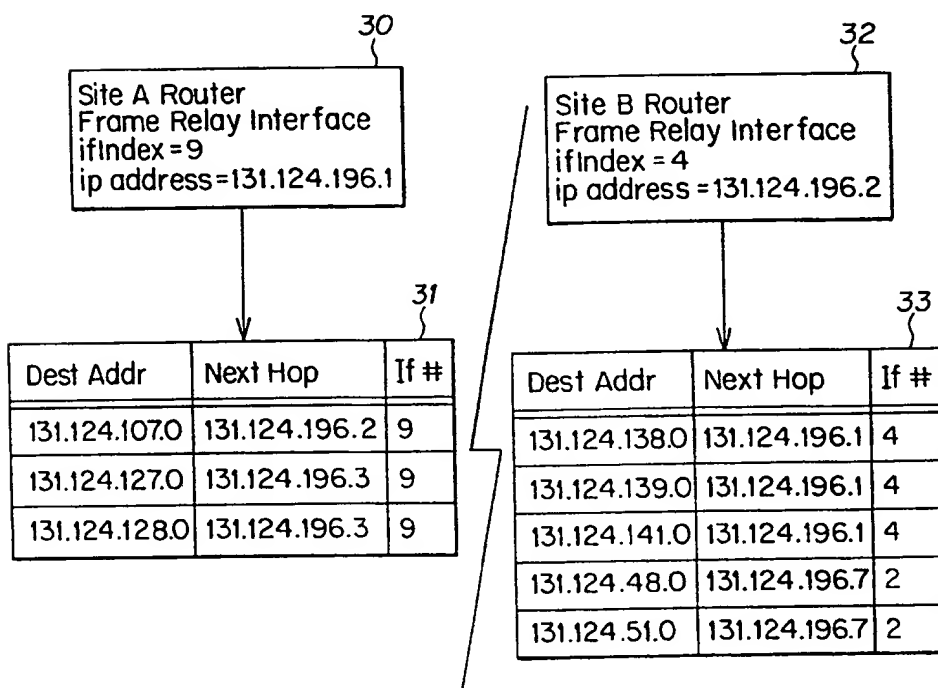


FIG. 3

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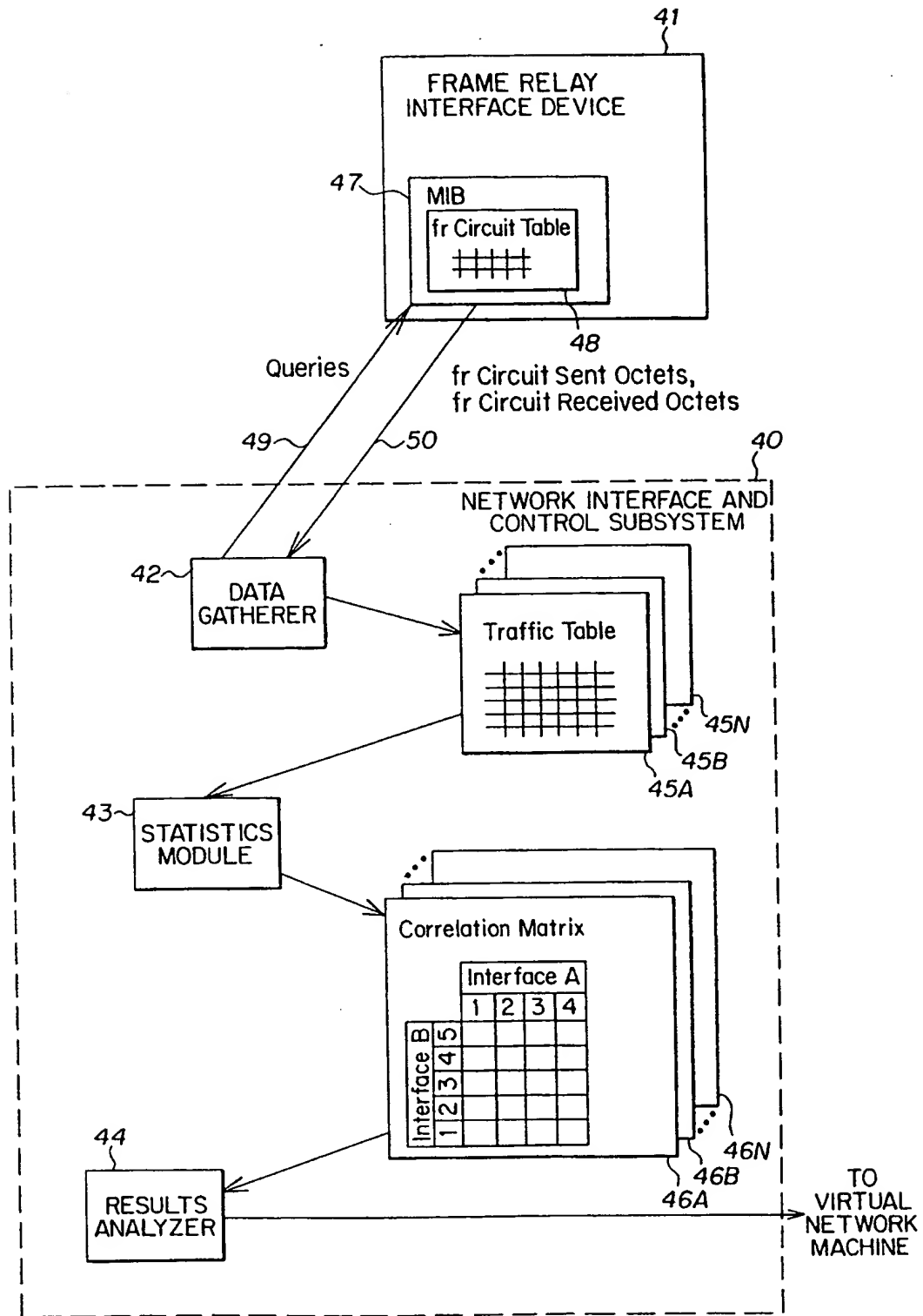
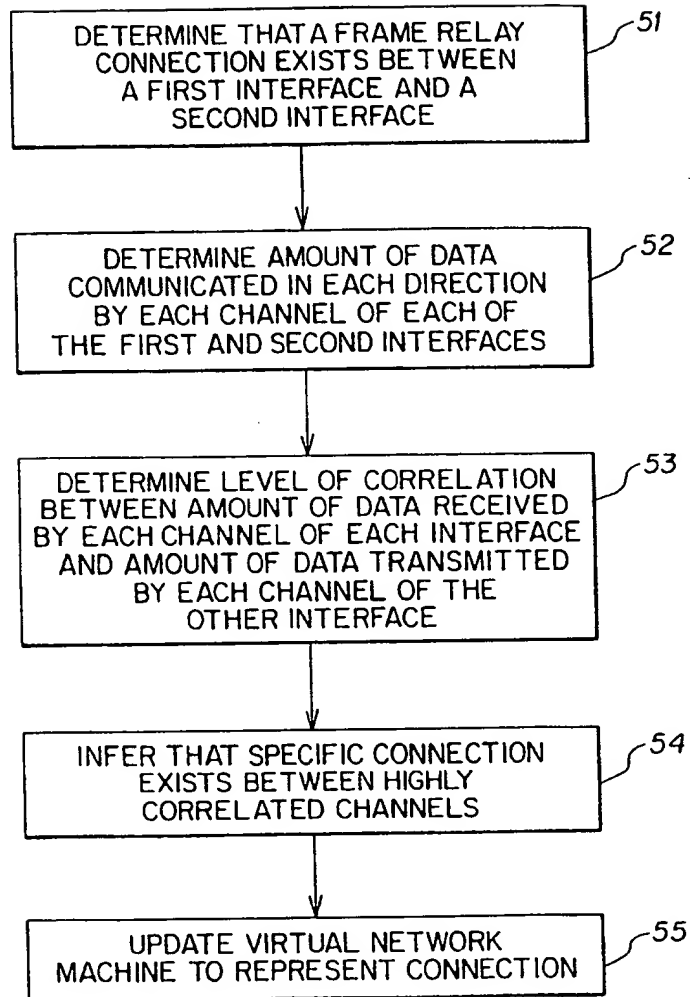
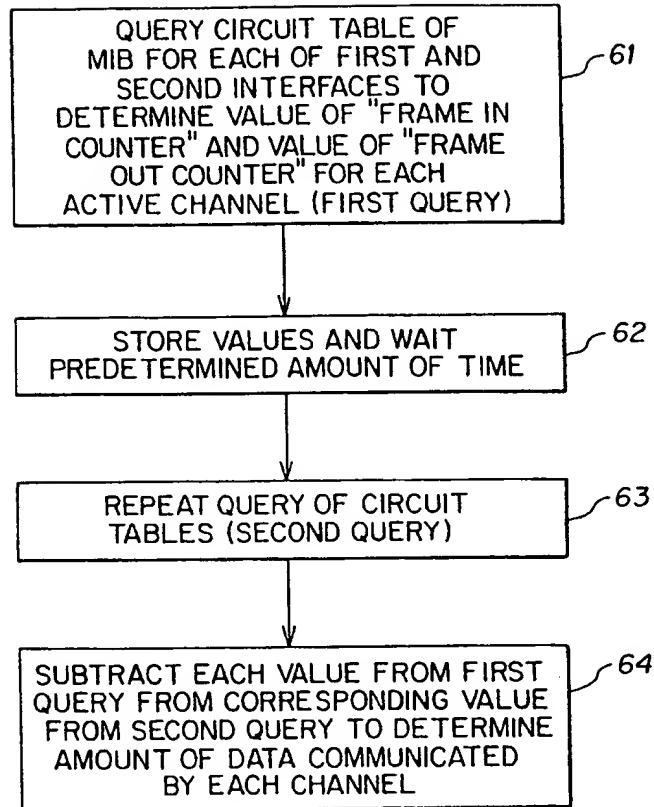


FIG. 4

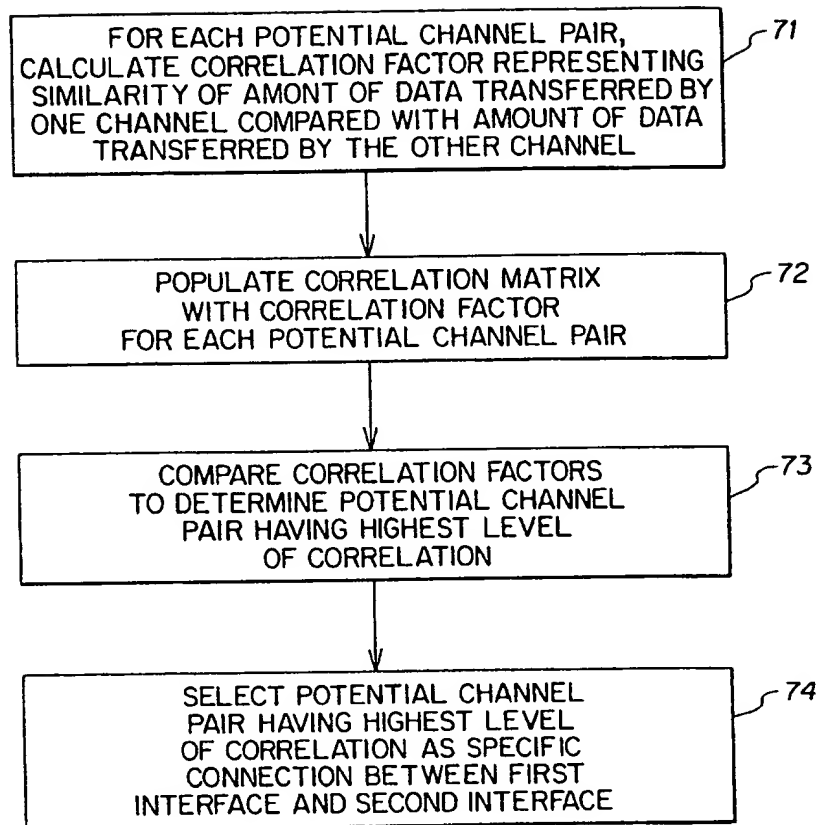
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**FIG. 5**

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**FIG. 6**

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**FIG. 7**

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81 IfIndex	82 DLCI	83 OutFrames-t0	84 InFrames-t0	85 OutFrames-t1	86 InFrames-t1	87 DeltaOut	88 DeltaIn	80
9	31	1,215,647	6,894,341	1,264,893	6,921,004	49,246	26,663	
9	103	32,423,675	37,654,687	32,498,011	37,721,956	74,336	67,269	
9	351	6,745,234	112,435	6,772,156	112,435	26,922	0	

FIG. 8

91 IfIndex	92 DLCI	93 OutFrames-t0	94 InFrames-t0	95 OutFrames-t1	96 InFrames-t1	97 DeltaOut	98 DeltaIn	90
4	40	14,342,552	8,553,341	14,356,689	8,554,067	14,137	726	
4	56	6,943,698	1,284,577	6,970,541	1,333,8232	26,843	49,246	

FIG. 9

		INTERFACE If 4		100
		DLCI 40	DLCI 56	
INTERFACE If 9	DLCI 31	32	99	
	DLCI 103	19	70	
	DLCI 351	4	52	

FIG. 10

INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 98/05370

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 H04L12/26

According to International Patent Classification(IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 6 H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	US 5 450 408 A (PHAAL) 12 September 1995 see abstract see column 2, line 1 - line 42 ---	1,9,17
A	EP 0 455 402 A (HEWLETT-PACKARD COMPANY) 6 November 1991 see abstract see column 2, line 31 - column 3, line 14 ---	1,9,17
A	WO 95 06989 A (CABLETRON SYSTEMS INC.) 9 March 1995 see abstract see page 9, line 35 - page 10, line 17 see page 2, line 28 - page 3, line 8 see page 3, line 14 - page 4, line 5 --- -/--	1,9,17

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

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- *T* later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention
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- *Y* document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.
- *&* document member of the same patent family

Date of the actual completion of the international search

26 August 1998

Date of mailing of the international search report

03/09/1998

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INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 98/05370

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	<p>MULLER N J: "THE MANAGEMENT ASPECTS OF FRAME RELAY" INTERNATIONAL JOURNAL OF NETWORK MANAGEMENT, vol. 2, no. 2, June 1992, pages 87-99, XP002052342 see page 91, right-hand column, line 20 - line 34 see page 94, left-hand column, line 46 - right-hand column, line 48 -----</p>	1-24

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